The number of connected clusters of total weight w = m + 1 that contain a is at most  $e \mathfrak{d} (1 + e(\mathfrak{d} - 1))^{w-1}$  by Proposition 136, hence we have Item (3). Since Proposition 138 gives the uniform bound

$$\left| \frac{1}{\mathbf{V}!} \mathcal{D}_{\mathbf{V}} \mathcal{L} \right| \leq \left( 2e(\mathfrak{d}+1) \beta \right)^{m+1},$$

and differentiating the monomial  $\lambda^{\mathbf{V}}$  contributes at most a factor (m+1) since  $|\partial_{\lambda_a}\lambda^{\mathbf{V}}| \leq \mu(a) \leq m+1$ . Dividing by  $\beta^{m+1}$  as in (47) yields Item (4): each monomial coefficient in  $p_m^{(a)}$  has size at most  $(2e(\mathfrak{d}+1))^{m+1}(m+1)$ .

Now let us briefly discuss Items (A) and (B) in Theorem 140. For Item (A), to enumerate all contributing monomials, one enumerates connected clusters of weight m rooted at a by a breadth-first, layer-by-layer procedure (see Algorithm 1, i.e. "tails" in [HKT22]). Given random-access to neighbors in  $\mathfrak{G}$ , the total time is  $O(m\mathfrak{d} C)$  where C is the number of clusters (hence monomials), giving Item (A).

For Item (B), to compute an individual coefficient exactly, [HKT22] shows how to evaluate the needed cluster derivatives  $\mathcal{D}_{\mathbf{V}}\mathcal{L}$  symbolically using faithful Pauli representations, in time  $O(Lm^3 + 8^m m^5 \log^2 m)$ .

## 3.2. Finding a solution using convexity

In the previous subsection we established a high-temperature expansion for the observables  $\langle E_a \rangle_{\beta} = \operatorname{tr}(E_a \rho_{\beta})$  and proved quantitative bounds on the size and locality of the resulting polynomials in Theorem 140. We now leverage those bounds to show that  $\mathcal{L}(\lambda) = \log \operatorname{tr}\left(e^{-\beta \sum_{a \in [M]} \lambda_a E_a}\right)$  is locally strongly convex in the high-temperature regime. This convexity will be the key ingredient that lets us robustly invert the map from Hamiltonian coefficients to thermal expectations, and thereby learn the coefficients.

Fix a vector  $x = (x_1, \dots, x_M) \in [-1, 1]^M$ . By Theorem 140 we may write

$$\langle E_a \rangle_{\beta}(x) = \sum_{m=1}^{\infty} \beta^m \, p_m^{(a)}(x) \,, \quad p_m^{(a)} \text{ homogeneous of degree } m,$$

where  $p_m^{(a)}$  only depends on entries  $x_b$  with  $\operatorname{dist}_{\mathfrak{G}}(a,b) \leq m$ , and its number and size of coefficients obey the bounds from Theorem 140 (Items (3)–(4)). In particular, letting

$$\tau := (1 + e(\mathfrak{d} - 1)) (2e(\mathfrak{d} + 1)) \le 2e^2(\mathfrak{d} + 1)^2 \tag{48}$$

as before, the sum of absolute coefficients of  $p_m^{(a)}$  is bounded by

$$c_m = e \,\mathfrak{d} \,(1 + e(\mathfrak{d} - 1))^m \,(2e(\mathfrak{d} + 1))^{m+1}(m+1)$$
  
=  $2e^2 \mathfrak{d}(\mathfrak{d} + 1) \,\tau^m(m+1)$ . (49)

For the learning task we will work with a *shifted, truncated* map  $\mathcal{F}: [-1,1]^M \to \mathbb{R}^M$  whose a-th component is

$$\mathcal{F}_a(x) := \sum_{m=0}^{m_{\text{max}}} \beta^m \, p_m^{(a)}(x) = -\widehat{E}_a - \beta \, x_a + \beta^2 p_2^{(a)}(x) + \dots + \beta^m p_{m_{\text{max}}}^{(a)}(x) \,, \, (50)$$

where  $\widehat{E}_a$  is an estimate of  $\langle E_a \rangle_{\beta}(\lambda)$  obtained from measurements (so we set  $p_0^{(a)} := -\widehat{E}_a$ ),  $p_1^{(a)}(x) = -x_a$  by a short computation, and  $m_{\text{max}}$  is a truncation order we

will later choose polylogarithmic in  $1/(\beta \varepsilon)$ . Our strategy will be to find an x such that  $\mathcal{F}_a(x)$  is small for all  $a \in [M]$ ; we will argue that if we can do then, then x is guaranteed to be closed to the true couplings  $\lambda$  by a convexity argument.

Here we will articulate our basic proof strategy. Let  $J(x) = d\mathcal{F}(x)$  be the Jacobian of  $\mathcal{F}$ , namely

$$J_{ab}(x) = \frac{\partial}{\partial x_b} \mathcal{F}_a(x) \,.$$

Recall that the norm  $\|\cdot\|_{\infty\to\infty}$  is defined by  $\|A\|_{\infty\to\infty} = \max_{x\neq 0} \frac{\|Ax\|_{\infty}}{\|x\|_{\infty}}$ . Then the idea is to use Newton iteration to find a x such that  $\|x-\lambda\|_{\infty} \leq O(\varepsilon)$ ; doing so will involve bounding the size of the inverse Jacobian  $J^{-1}$  which plays an important role in Newton iteration, as well as the size of  $\mathcal{F}(\lambda)$  which is the target value of  $\mathcal{F}$ .

To prepare for our Newton's method procedure, we will want to first establish the following facts:

- (1) For suitable conditions on  $\beta$  and  $\mathfrak{d}$ , we have  $||J(x)^{-1}||_{\infty\to\infty} \leq 2\beta^{-1}$  for all  $m_{\max} \geq 1$ .
- (2) For any  $\varepsilon > 0$ , we can choose  $m_{\text{max}}$  sufficiently large (with suitable conditions on  $\beta$  and  $\mathfrak{d}$ ) such that  $\|\mathcal{F}(\lambda)\|_{\infty} \leq O(\beta \varepsilon)$ .

For the first condition, we really only need the condition to hold for  $m_{\text{max}}$  sufficiently large, but in fact we will show that it holds for all  $m_{\text{max}} \geq 1$ .

We will begin by establishing the first condition, and then treat the second. To this end, we have the following lemma.

## Lemma 141. Suppose that

$$100e^{6}(\mathfrak{d}+1)^{8}\beta \le 1. {(51)}$$

Then for any  $x \in [-1,1]^M$ , we have  $\|\mathbb{1} + \beta^{-1}J(x)\|_{\infty \to \infty} \le \frac{1}{2}$  and  $\|J(x)^{-1}\|_{\infty \to \infty} \le 2\beta^{-1}$  for any  $m_{\max} \ge 1$ .

PROOF. We note that if  $\|\mathbb{1} + \beta^{-1}J\|_{\infty \to \infty} \leq \frac{1}{2}$ , then since

$$J^{-1} = -\frac{1}{\beta} \frac{1}{1 - (1 - \beta^{-1}J)} = -\frac{1}{\beta} \sum_{k=0}^{\infty} (1 + \beta^{-1}J)^k,$$

we would have

$$||J^{-1}||_{\infty \to \infty} \le \beta^{-1} \sum_{k=0}^{\infty} ||\mathbb{1} + \beta^{-1}J||_{\infty \to \infty}^k \le 2\beta^{-1}.$$

Thus it suffices to show  $\|\mathbb{1}+\beta^{-1}J(x)\|_{\infty\to\infty} \leq \frac{1}{2}$ , or equivalently  $\|\beta\mathbb{1}+J(x)\|_{\infty\to\infty} \leq \frac{\beta}{2}$ , for our stated domain of  $\beta$ .

We observe that the Jacobian takes the form

$$J_{ab} = -\beta \, \delta_{ab} + O(\beta^2) \,,$$

and so  $\beta \mathbb{1} + J = O(\beta^2)$ . As such, we would like to bound the  $O(\beta^2)$  remainder. Let  $u = (u_1, ..., u_M)$  satisfy  $|u_b| \leq 1$  for all b, i.e.  $||u||_{\infty} \leq 1$ . Then we have

$$((J + \beta \mathbb{1})u)_a = \sum_b (J + \beta \mathbb{1})_{ab} u_b$$

$$= \sum_b u_b \left( \beta^2 \frac{\partial}{\partial x_b} p_2^{(a)}(x) + \dots + \beta^{m_{\text{max}}} \frac{\partial}{\partial x_b} p_{m_{\text{max}}}^{(a)}(x) \right)$$

$$= \sum_{k=2}^{m_{\text{max}}} \beta^k \sum_{b: \text{dist}_{\mathcal{A}}(a,b) \le k} u_b \frac{\partial}{\partial x_b} p_k^{(a)}(x) ,$$

where in going to the last line we have used Item (2) of Theorem 140. We observe that in the last sum, at each fixed k, the index b ranges over at most  $1+\mathfrak{d}+\cdots+\mathfrak{d}^k \leq (\mathfrak{d}+1)^k$  vertices of  $\mathfrak{G}$ . Now recall that each  $p_k^{(a)}$  is a homogeneous polynomial of degree k, and that the sum of the absolute values of the coefficients is bounded by  $c_k$  in (49). Therefore  $\left|\frac{\partial}{\partial x_b}p_k^{(a)}\right| \leq kc_k$  in the domain of  $\mathcal{F}$ , and we have

$$\begin{split} |((J+\beta \mathbb{1})u)_a| & \leq \sum_{k=2}^{\infty} \beta^k (\mathfrak{d}+1)^k k c_k \\ & \leq 2 e^2 (\mathfrak{d}+1)^2 (\beta (\mathfrak{d}+1)\tau)^2 \sum_{k=2}^{\infty} (\beta (\mathfrak{d}+1)\tau)^{k-2} k (k+1) \\ & = 2 e^2 (\mathfrak{d}+1)^4 \beta^2 \tau^2 \left( \left. \frac{6-6r+2r^2}{(1-r^3)} \right|_{r=\beta (\mathfrak{d}+1)\tau} \right) \\ & \leq \frac{25}{2} \, e^2 (\mathfrak{d}+1)^4 \beta^2 \tau^2 \, . \end{split}$$

In going from the second line to the third line we used  $\beta(\mathfrak{d}+1)\tau < 1$ , and in going to the last line we used  $\beta(\mathfrak{d}+1)\tau \leq \frac{1}{100}$ . Since our u satisfying  $||u||_{\infty} \leq 1$  was arbitrary, we have obtained the bound  $||J+\beta\mathbb{1}||_{\infty\to\infty} \leq \frac{25}{2}\,e^2(\mathfrak{d}+1)^4\beta^2\tau^2$ . Using  $\tau \leq 2e^2(\mathfrak{d}+1)^2$  from (48) and  $100e^6(\mathfrak{d}+1)^8\beta \leq 1$  from (51), we find our desired bound  $||J+\beta\mathbb{1}||_{\infty\to\infty} \leq \frac{\beta}{2}$ .

A nice consequence of the above lemma is the following convexity result:

**Lemma 142.** If (51) holds, then  $\nabla^{\otimes 2} \mathcal{L} \succeq \frac{\beta^2}{2} \mathbb{1}$ , namely  $\mathcal{L}$  is  $(\frac{\beta^2}{2})$ -strongly convex.

PROOF. Take  $m_{\max} = \infty$  so that  $\nabla^{\otimes 2} \mathcal{L} = -\beta J$ , where we note that the Jacobian J is Hermitian. For a Hermitian matrix X, we have  $||X|| \leq ||X||_{\infty \to \infty}$ , and so  $||\mathbb{1} + \beta^{-1}J|| \leq ||\mathbb{1} + \beta^{-1}J||_{\infty \to \infty} \leq \frac{1}{2}$ , implying that  $\mathbb{1} + \beta^{-1}J \leq \mathbb{1}/2$  and thus  $\beta^{-1}J \leq -\mathbb{1}/2$ , which is equivalent to  $\nabla^{\otimes 2} \mathcal{L} \succeq \frac{\beta^2}{2} \mathbb{1}$ .

Next let us show that if  $m_{\text{max}}$  is chosen to scale at least logarithmically in  $1/(\beta \varepsilon)$ , then we can arrange for  $\|\mathcal{F}(\lambda)\|_{\infty} \leq O(\beta \varepsilon)$ . First we require the following lemma.

**Lemma 143** (Estimating thermal expectations in parallel). For any  $\varepsilon, \delta \in (0,1)$  there is a measurement procedure that (given independent copies of  $\rho_{\beta}$ ) produces estimators  $\widehat{E}_a$  such that

$$\left|\widehat{E}_a - \langle E_a \rangle_\beta\right| \le \beta \, \varepsilon \quad \text{for all } a \in [M]$$

simultaneously with probability at least  $1 - \delta$ , using

$$O\left(\frac{\mathfrak{d}}{\beta^2 \varepsilon^2} \log \frac{M}{\delta}\right)$$

copies of  $\rho_{\beta}$  and with time complexity

$$O\left(\frac{N\mathfrak{d}}{\beta^2\varepsilon^2}\log\frac{M}{\delta}\right).$$

PROOF. We first recall a standard fact: given a quantum state  $\rho$  and a Hermitian observable E with  $||E|| \leq 1$ , one can estimate  $\operatorname{tr}(E\rho)$  to additive error  $\varepsilon_0$  with success probability at least  $1 - \delta_0$  using  $O(\log(1/\delta_0)/\varepsilon_0^2)$  independent copies of  $\rho$ . Indeed, measuring  $\rho$  in the eigenbasis of E yields an i.i.d. random variable in [-1,1] whose expectation is  $\operatorname{tr}(E\rho)$ ; Hoeffding bounds then give the stated sample complexity.

We now apply this in parallel to the family  $\{E_a\}_{a\in[M]}$ . Color the vertices of the dual interaction graph  $\mathfrak G$  using at most  $\mathfrak d+1$  colors (a greedy coloring suffices). By definition of  $\mathfrak G$ , all  $E_a$  belonging to a fixed color class act on disjoint sets of qubits. Consequently, on a single copy of  $\rho_\beta$  we can measure all  $E_a$  in that color class simultaneously: since each  $E_a$  is a Pauli string, it suffices to measure each qubit once in the appropriate single-qubit Pauli basis and multiply outcomes to obtain the eigenvalue of each  $E_a$  in the class.

Fix a color class and set the target accuracy per observable to  $\varepsilon_0 := \beta \varepsilon$ . By the single-observable estimate and a union bound over all a in the class,  $O\left(\log(1/\delta_0)/\varepsilon_0^2\right)$  copies of  $\rho_\beta$  suffice to ensure that every  $\widehat{E}_a$  in that class satisfies  $|\widehat{E}_a - \langle E_a \rangle_\beta| \le \varepsilon_0$  with probability at least  $1 - \delta_0$ . Repeating independently for each of the at most  $\mathfrak{d} + 1$  color classes, the total number of copies is

$$(\mathfrak{d}+1)\,O\Big(\frac{\log(1/\delta_0)}{\varepsilon_0^2}\Big) = O\bigg(\frac{\mathfrak{d}}{\beta^2\varepsilon^2}\,\log\frac{1}{\delta_0}\bigg)\,.$$

Choosing  $\delta_0 := \delta/M$  and applying a union bound across all M observables yields simultaneous accuracy  $|\widehat{E}_a - \langle E_a \rangle_{\beta}| \leq \beta \varepsilon$  for every  $a \in [M]$  with probability at least  $1 - \delta$ , and the stated copy complexity  $O\left(\frac{\mathfrak{d}}{\beta^2 \varepsilon^2} \log \frac{M}{\delta}\right)$  follows.

For the time complexity, note that each copy used in a given color round requires at most N single-qubit Pauli measurements (one per qubit), and there are  $(\mathfrak{d} + 1) O(\log(1/\delta_0)/\varepsilon_0^2)$  such copies overall. This gives time

$$O\left(N\frac{\mathfrak{d}}{\beta^2\varepsilon^2}\log\frac{M}{\delta}\right),$$

as claimed.  $\Box$ 

This lemma tells us that we can set  $|\widehat{E}_a - \langle E_a \rangle_{\beta}| \leq O(\beta \varepsilon)$  for all a. With this in mind, we have the following.

**Lemma 144.** Assume the high-temperature condition (51). Let  $\tau$  be as in (48) and set  $r := \beta \tau$ . Suppose the empirical means obey  $|\widehat{E}_a - \langle E_a \rangle_{\beta}| \leq \beta \varepsilon$  for all  $a \in [M]$ . If the truncation order  $m_{\max}$  in (50) satisfies

$$(2r)^{m_{\max}+1} \le \frac{\beta \varepsilon}{4 e^2 \mathfrak{d}(\mathfrak{d}+1)}, \tag{52}$$

then  $\|\mathcal{F}(\lambda)\|_{\infty} \leq 2\beta\varepsilon$ . Equivalently, it suffices to take

$$m_{\text{max}} \ge \left\lceil \frac{\log\left(\frac{4e^2\mathfrak{d}(\mathfrak{d}+1)}{\beta\varepsilon}\right)}{\log\left(\frac{1}{2\beta\tau}\right)} \right\rceil - 1.$$
 (53)

In particular, for constant  $\mathfrak{d}$  we have  $m_{\text{max}} = O(\log(1/(\beta \epsilon)))$ .

PROOF. By the definition (50) and the triangle inequality,

$$\left| \mathcal{F}_a(\lambda) \right| \le \left| \widehat{E}_a - \langle E_a \rangle_{\beta} \right| + \sum_{m > m_{\max}} \beta^m \left| p_m^{(a)}(\lambda) \right| \le \beta \varepsilon + \sum_{m > m_{\max}} \beta^m c_m.$$

Thus, with  $r = \beta \tau$ ,

$$\sum_{m>m_{\max}} \beta^m c_m = 2e^2 \mathfrak{d}(\mathfrak{d}+1) \sum_{m>m_{\max}} (m+1) r^m.$$

For  $m \ge 1$  we may use  $(m+1) \le 2^m$ , where (since 2r < 1)

$$\sum_{m > m_{\text{max}}} (m+1) r^m \le \sum_{m > m_{\text{max}}} (2r)^m = \frac{(2r)^{m_{\text{max}}+1}}{1-2r}.$$

The high-temperature hypothesis (51) implies  $r \ll 1$  (and hence 2r < 1); in particular,  $1/(1-2r) \le 2$ . Therefore

$$\sum_{m>m_{\max}} \beta^m c_m \le 4e^2 \mathfrak{d}(\mathfrak{d}+1) (2r)^{m_{\max}+1}.$$

Imposing (52) makes the right-hand side at most  $\beta \varepsilon$ , and hence  $|\mathcal{F}_a(\lambda)| \leq 2\beta \varepsilon$  for all a. Taking the maximum over a yields  $||\mathcal{F}(\lambda)||_{\infty} \leq 2\beta \varepsilon$ .

Finally, solving (52) for  $m_{\text{max}}$  gives (53); since  $2\beta\tau < 1$  under (51), the denominator is a positive constant when  $\mathfrak{d}$  is constant, proving the claimed  $O(\log(1/(\beta\varepsilon)))$  scaling.

Finally, we show that we can efficiently find an x such that  $||x - \lambda||_{\infty} \le 18 \varepsilon$ .

Theorem 145 (High-temperature learning via projected Newton-Raphson). Assume the high-temperature condition (51). Suppose we are given estimates  $\{\widehat{E}_a\}_{a\in[M]}$  of the thermal expectations  $\langle E_a\rangle_{\beta}$  obeying  $|\widehat{E}_a - \langle E_a\rangle_{\beta}| \leq \beta \varepsilon$  for all  $a \in [M]$ . Moreover let us take  $\varepsilon \leq \frac{1}{12}$ . Then there is a classical algorithm (a projected Newton-Raphson scheme with a truncated Neumann-series inverse) that outputs  $x \in [-1,1]^M$  such that  $\|x-\lambda\|_{\infty} \leq 18\varepsilon$  in time  $O(\frac{ML}{\varepsilon}\operatorname{poly}(\mathfrak{d},\log\frac{1}{\beta\varepsilon}))$ , where L is the maximum number of qubits on which any Hamiltonian term acts.

PROOF SKETCH. Let us choose the judicious bound

$$m_{\max} \ge \left\lceil \frac{e}{e-1} \frac{1}{\log\left(\frac{1}{\beta\tau}\right)} \log\left(\frac{12e^2(\mathfrak{d}+1)^2}{\beta\varepsilon\log\left(\frac{1}{\beta\tau}\right)}\right) \right\rceil$$

which is compatible with our previous one. The **Newton-Raphson method** ordinarily entails an iteration like  $x^{(t+1)} = x^{(t)} - (J^{-1}\mathcal{F})(x^{(t)})$ , although to avoid computing the inverse of J we will instead consider an approximation  $J(x)^{-1} \approx$ 

 $\beta^{-1} \sum_{k=0}^{K-1} (\mathbb{1} + \beta^{-1} J(x))^k$  for a sufficiently large K that we will specify. Specifically, we consider the iteration

$$x^{(0)} = \vec{0}, \quad x^{(t+1)} = \operatorname{Proj}_{[-1,1]^M} \left[ x^{(t)} + \beta^{-1} \sum_{k=1}^K (\mathbb{1} + \beta^{-1} J(x^{(t)}))^k \mathcal{F}(x^{(t)}) \right]$$

where we have used

$$\operatorname{Proj}(u) := \begin{cases} 1 & \text{if } u \in (1, \infty) \\ u & \text{if } u \in [-1, 1] \\ -1 & \text{if } u \in (-\infty, 1) \end{cases},$$

and take  $K = \lceil \log_2(\frac{3}{2\varepsilon}) \rceil$ .

Before analyzing the convergence of the iterations, let us examine the error  $e^{(t)}$  between  $J(x^{(t)})^{-1}\mathcal{F}(x^{(t)})$  and  $\beta^{-1}\sum_{k=0}^{K-1}(\mathbb{1}+\beta^{-1}J(x^{(t)}))^k\mathcal{F}(x^{(t)})$ . Specifically, we have

$$e^{(t)} := \left( -J(x^{(t)})^{-1} + \frac{1}{\beta} \sum_{k=0}^{K-1} (\mathbb{1} + \beta^{-1} J(x^{(t)}))^k \right) \mathcal{F}(x^{(t)})$$

$$= -\frac{1}{\beta} \sum_{k=K}^{\infty} (\mathbb{1} + \beta^{-1} J(x^{(t)})^k \mathcal{F}(x^{(t)})$$

$$= J^{-1}(x^{(t)}) (\mathbb{1} + \beta^{-1} J(x^{(t)}))^K \mathcal{F}(x^{(t)}),$$

which by Lemma 141 decays exponentially in K in the  $\|\cdot\|_{\infty}$  norm. This will be useful for us shortly.

With the error  $e^{(t)}$  under control, let us examine the convergence of  $x^{(t)}$  under our Newton-Raphson iteration. Let  $\mathcal{F}_a(s):[0,1]\to\mathbb{R}$  by defined by  $\mathcal{F}_a(s):=\mathcal{F}_a(x+s(\lambda-x))$ . By Taylor's remainder theorem, there exists an  $s'\in[0,1]$  such that

$$\underbrace{\mathcal{F}_a(1)}_{=\mathcal{F}_a(\lambda)} = \underbrace{\mathcal{F}_a(0)}_{=\mathcal{F}_a(x)} + (\partial_s \mathcal{F}_a)(0) + \frac{1}{2} (\partial_s^2 \mathcal{F}_a)(s').$$

Using  $\partial_s = \sum_b (\lambda_b - x_b) \partial_b$  and setting  $y^{(a)} := s'\lambda + (1 - s')x$ , we find

$$\mathcal{F}_a(\lambda) = \mathcal{F}_a(x) + \sum_b (\lambda_b - x_b) \underbrace{(\partial_b \mathcal{F}_a)(x)}_{=J_{ab}(x)} + \frac{1}{2} \sum_{b,c} (\lambda_b - x_b)(\lambda_c - x_c)(\partial_b \partial_c \mathcal{F}_a)(y^{(a)}).$$

Letting  $\Delta^{(t)} := x^{(t)} - \lambda$  (and similarly  $\Delta^{(t+1)} := x^{(t+1)} - \lambda$ ) where  $\Delta_d^{(t)}$  denotes the dth coordinate, we have the following:

$$\begin{split} |\Delta_{d}^{(t+1)}| &= \left| \text{Proj}_{[-1,1]}[(x - (J^{-1}\mathcal{F})(x) + e)_{d}] - \lambda_{d} \right| \\ &\leq \left| (x - (J^{-1}\mathcal{F})(x) + e)_{d} - \lambda_{d} \right| \\ &= \left| e_{d}^{(t)} + \Delta_{d}^{(t)} - \sum_{a} (J(x^{(t)})^{-1})_{da} \mathcal{F}_{a}(x^{(t)}) \right| \\ &= \left| e_{d}^{(t)} + \Delta_{d}^{(t)} - \sum_{a} J(x^{(t)})_{da}^{-1} \left( \mathcal{F}_{a}(\lambda) - \sum_{b} (\lambda_{b} - x_{b}^{(t)}) J_{ab}(x^{(t)}) \right) - \frac{1}{2} \sum_{b,c} (\lambda_{b} - x_{b}^{(t)}) (\lambda_{c} - x_{c}^{(t)}) [\partial_{b} \partial_{c} \mathcal{F}_{a}](y^{(a)}) \right) \right| \\ &= \left| \left[ e^{(t)} + \Delta^{(t)} - J(x^{(t)})^{-1} \mathcal{F}(\lambda) - J(x^{(t)})^{-1} J(x^{(t)}) \Delta^{(t)} \right]_{d} + \frac{1}{2} \sum_{a,b,c} J(x^{(t)})_{da}^{-1} \Delta_{b}^{(t)} \Delta_{c}^{(t)} [\partial_{b} \partial_{c} \mathcal{F}](y^{(a)}) \right| \\ &= \left| \left[ J(x^{(t)})^{-1} \left( (\mathbb{1} + \beta^{-1} J(x^{(t)}))^{K} \mathcal{F}(x^{(t)}) - \mathcal{F}(\lambda) \right) \right]_{d} + \frac{1}{2} \sum_{a,b,c} J(x^{(t)})_{da}^{-1} \Delta_{b}^{(t)} \Delta_{c}^{(t)} [\partial_{b} \partial_{c} \mathcal{F}_{a}](y^{(a)}) \right| . (54) \end{split}$$

We will bound each term in the last equation in turn. For the first part, we have

$$\left| \left[ J(x^{(t)})^{-1} \left( (\mathbb{1} + \beta^{-1} J(x^{(t)}))^K \mathcal{F}(x^{(t)}) - \mathcal{F}(\lambda) \right) \right]_d \right|$$

$$\leq \|J(x^{(t)})^{-1}\|_{\infty \to \infty} \left( \|\mathbb{1} + \beta^{-1} J(x^{(t)})\|_{\infty \to \infty}^K \|\mathcal{F}(x^{(t)})\|_{\infty} + \|\mathcal{F}(\lambda)\|_{\infty} \right)$$

$$\leq 2\beta^{-1} \left( 2^{-K} (2 + \beta \varepsilon) + 2\beta \varepsilon \right) \leq 6\varepsilon .$$
(55)

In going to the last line we have used Lemma 141 and Lemma 144, in conjunction with

$$|\mathcal{F}_{a}(x)| \leq \left| \widehat{E}_{a} + \sum_{k=1}^{m_{\text{max}}} \beta^{k} |p_{k}^{(a)}(x)| \right|$$

$$\leq |\widehat{E}_{a} - \langle E_{a} \rangle_{\beta}| + \left| -\langle E_{a} \rangle_{\beta} + \sum_{k=1}^{m_{\text{max}}} \beta^{k} |p_{k}^{(a)}(x)| \right|$$

$$\leq \beta \varepsilon + 2.$$

For the last term in (54), we have for all indices d the inequalities

$$\left| \frac{1}{2} \sum_{a,b,c} J(x^{(t)})_{da}^{-1} \Delta_{b}^{(t)} \Delta_{c}^{(t)} [\partial_{b} \partial_{c} \mathcal{F}](y^{(a)}) \right|$$

$$\leq \frac{1}{2} \|J(x^{(t)})^{-1}\|_{\infty \to \infty} \max_{a} \left| \sum_{b,c} \Delta_{b}^{(t)} \Delta_{c}^{(t)} [\partial_{b} \partial_{c} \mathcal{F}_{a}](y^{(a)}) \right|$$

$$\leq \frac{1}{\beta} \max_{a} \sum_{k=0}^{\infty} \sum_{b,c} |\Delta_{b}^{(t)} \Delta_{c}^{(t)}| \beta^{k} |\partial_{b} \partial_{c} p_{k}^{(a)}(y)|$$

$$\leq \frac{1}{\beta} \sum_{k=0}^{\infty} \sum_{\substack{b,c:\\ \text{dist}_{\mathfrak{G}}(b,a) \leq k\\ \text{dist}_{\mathfrak{G}}(c,a) \leq k}} \|\Delta^{(t)}\|_{\infty}^{2} \beta^{k} k(k-1) c_{k}$$

$$\leq \frac{1}{\beta} \sum_{k=0}^{\infty} (\mathfrak{d}+1)^{2k} \|\Delta^{(t)}\|_{\infty}^{2} \beta^{k} k(k-1) c_{k}$$

$$= \frac{12e^{2}}{\beta} \|\Delta^{(t)}\|_{\infty}^{2} (\mathfrak{d}+1)^{2} \frac{(\beta(\mathfrak{d}+1)^{2}\tau)^{2}}{(1-\beta(\mathfrak{d}+1)^{2}\tau)^{4}}$$

$$\leq \frac{25}{2} e^{2} \beta(\mathfrak{d}+1)^{6} \tau^{2} \|\Delta^{(t)}\|_{\infty}^{2}, \qquad (56)$$

where in going to the second-to-last line we have used that  $\beta(\mathfrak{d}+1)^2\tau < 1$  and in going to the last line we have used that  $\beta\mathfrak{d}^2\tau \leq 1 - \left(\frac{24}{25}\right)^{1/4}$ . Putting together (55) and (56) we find

$$\|\Delta^{(t+1)}\|_{\infty} \le 6\varepsilon + \frac{25}{2} e^2 \beta (\mathfrak{d} + 1)^6 \tau^2 \|\Delta^{(t)}\|_{\infty}^2.$$

By solving the recursion, one can show that so long as  $\|\Delta^{(0)}\|_{\infty} \leq \frac{1}{25e^2\beta(\mathfrak{d}+1)^6\tau^2} \leq 1$ , we achieve  $\|x^{(T)} - \lambda\|_{\infty} \leq 18\varepsilon$  for

$$T = \lceil -\log_2(300e^6(\mathfrak{d}+1)^{10}\beta\varepsilon) \rceil.$$

Finally, let us sketch the runtime bound. For each  $a \in [M]$ , the truncated series  $\mathcal{F}_a(x) = \sum_{k=0}^m \beta^k p_k^{(a)}(x) - \widehat{E}_a$  is a degree-m polynomial whose support is contained in the radius-k neighborhoods of a in  $\mathfrak{G}$ ; the number of contributing terms at order k is at most poly( $\mathfrak{d}$ ) ( $\mathfrak{d}+1$ )<sup>k</sup> and each coefficient can be evaluated in time  $O(L \operatorname{poly}(k))$ . Hence, evaluating all M coordinates of F(x) and forming (or applying) the nonzeros of the sparse Jacobian J(x) at a given point x costs

$$O\Big(M L \operatorname{poly}(\mathfrak{d}) \sum_{k=0}^{m} (\mathfrak{d}+1)^k\Big) = O\Big(M L \operatorname{poly}(\mathfrak{d}) (\mathfrak{d}+1)^{O(m)}\Big).$$

One Newton step uses the truncated Neumann-series inverse  $\beta^{-1} \sum_{k=0}^{K-1} (\mathbb{1} + \beta^{-1} J(x))^k$ , which requires K sparse matrix-vector multiplies with J(x), and thus has cost  $O(K M L \operatorname{poly}(\mathfrak{d}) (\mathfrak{d} + 1)^{O(m)})$  at iteration  $x = x^{(t)}$ . The projection  $\operatorname{Proj}_{[-1,1]^M}$  adds only O(M) time. With T Newton iterations in total, the overall runtime is

$$O((K+1)TM L \operatorname{poly}(\mathfrak{d}) (\mathfrak{d}+1)^{O(m)}).$$

Substituting in our parameter choices yields the stated time complexity:

$$O\left(\frac{ML}{\varepsilon}\operatorname{poly}\left(\mathfrak{d},\log\frac{1}{\beta\varepsilon}\right)\right).$$

To summarize, we have succeeded in establishing that, for suitable  $\beta$ ,  $\mathfrak{d}$ , and  $m_{\max}$ , we have  $\|x - \lambda\|_{\infty} \leq O(\varepsilon)$ . Below we will put together all of our results thus far to get the final, overarching algorithm and associated bounds.

## 3.3. Putting the bounds together

We can combine all of the results above to get the main result of [HKT22]. We recapitulate some of the notation we have collected along the way.

Theorem 146 (Learning from high-temperature Gibbs states). Let  $H = \sum_{a \in [M]} \lambda_a E_a$  be a low-intersection Hamiltonian on N qubits: each non-identity Pauli term  $E_a$  acts on at most L = O(1) qubits and the dual interaction graph has maximum degree  $\mathfrak{d} = O(1)$ . Fix inverse temperature  $\beta > 0$  obeying the high-temperature condition (51) (equivalently  $\beta < \beta_c(\mathfrak{d})$  for a universal constant  $\beta_c > 0$  depending only on  $\mathfrak{d}$ ), and let  $\rho_\beta = e^{-\beta H}/\mathrm{tr}(e^{-\beta H})$ .

For any  $\varepsilon \in (0, \frac{1}{12})$  and failure probability  $\delta \in (0, 1)$ , there is a classical algorithm which, given independent copies of  $\rho_{\beta}$ , outputs  $\hat{\lambda} \in [-1, 1]^M$  satisfying

$$\|\widehat{\lambda} - \lambda\|_{\infty} \le 18\varepsilon$$

with probability at least  $1 - \delta$ , using

$$S_{\infty} = O\left(\frac{\mathfrak{d}}{\beta^2 \varepsilon^2} \log \frac{M}{\delta}\right)$$

copies of  $\rho_{\beta}$ . In particular, when  $\mathfrak{d} = O(1)$  and  $M = \Theta(N)$ , this is

$$S_{\infty} = O\left(\frac{\log N}{\beta^2 \varepsilon^2}\right).$$

Consequently, to achieve  $\ell_2$ -error  $\|\widehat{\lambda} - \lambda\|_2 \le \varepsilon$  it suffices to use

$$S_2 = O\left(\frac{M}{\beta^2 \varepsilon^2} \, \log \frac{M}{\delta}\right) = O\left(\frac{N}{\beta^2 \varepsilon^2} \, \log \frac{N}{\delta}\right).$$

The total running time is linear in the sample size (i.e. O(SN) where S is the number of copies used), up to polylogarithmic factors in  $1/(\beta \varepsilon)$ .

PROOF. Assume (51) and let  $\tau$  be as in (48). We can estimate all thermal expectations in parallel (via Lemma 143) to obtain  $\{\widehat{E}_a\}_{a\in[M]}$  with  $|\widehat{E}_a-\langle E_a\rangle_{\beta}| \leq \beta\varepsilon$  for every a using  $S_{\infty} = O(\frac{\mathfrak{d}}{\beta^2 z^2}\log\frac{M}{\lambda})$  copies of  $\rho_{\beta}$ , with success probability  $\geq 1 - \delta$ .

Define  $\mathcal{F}$  as in (50) and choose the truncation order  $m_{\text{max}}$  as in (53). Then Lemma 144 gives  $\|\mathcal{F}(\lambda)\|_{\infty} \leq 2\beta\varepsilon$ . By the high-temperature conditioning in Lemma 141, we have

$$\|\mathbb{1} + \beta^{-1} J(x)\|_{\infty \to \infty} \le \frac{1}{2}$$
 and  $\|J(x)^{-1}\|_{\infty \to \infty} \le 2\beta^{-1}$ , for all  $x \in [-1, 1]^M$ .

We run the projected Newton–Raphson update with truncated Neumann inverse in Theorem 145 from  $x^{(0)} = \vec{0}$  and with  $K = \lceil \log_2 \left(\frac{3}{2\varepsilon}\right) \rceil$ . The one-step analysis yields the recursion  $\|\Delta^{(t+1)}\|_{\infty} \le 6\varepsilon + C\beta \|\Delta^{(t)}\|_{\infty}^2$  with  $C = \frac{25}{2}e^2(\mathfrak{d}+1)^6\tau^2$ .

Solving this recursion with  $T = \left\lceil -\log_2\left(300e^6(\mathfrak{d}+1)^{10}\beta\varepsilon\right)\right\rceil$  gives  $\|x^{(T)}-\lambda\|_{\infty} \leq 18\varepsilon$ . We set  $\widehat{\lambda} := x^{(T)}$  to obtain the claimed accuracy with probability  $\geq 1 - \delta$ .

The sample bound is exactly that of Lemma 143, and for  $\mathfrak{d} = O(1)$  and  $M = \Theta(N)$  it simplifies to  $S_{\infty} = O(\frac{\log N}{\beta^2 \varepsilon^2})$ . The  $\ell_2$  statement follows by targeting  $\|\widehat{\lambda} - \lambda\|_{\infty} \le \varepsilon / \sqrt{M}$ , which replaces  $\varepsilon$  by  $\varepsilon / \sqrt{M}$  in Lemma 143, yielding  $S_2 = O(\frac{M}{\beta^2 \varepsilon^2} \log \frac{M}{\delta})$ . The runtime is  $O(S_{\infty}N)$  for data collection plus  $O(\frac{ML}{\varepsilon} \operatorname{poly}(\mathfrak{d}, \log \frac{1}{\beta \varepsilon}))$  for classical postprocessing, which is linear in the sample size up to polylogarithmic factors in  $1/(\beta \varepsilon)$ .